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#### Abstract

Various algorithm such as Doolittle, Crouts and Choleskyŝs have been proposed to factor a square matrix into a product of $L$ and $U$ matrices, that is, to find $L$ and $U$ such that $A=L U$; where $L$ and $U$ are lower and upper triangular matrices respectively. These methods are derived by writing the general forms of $L$ and $U$ and the unknown elements of $L$ and $U$ are then formed by equating the corresponding entries in $A$ and $L U$ in a systematic way. This approach for computing $L$ and $U$ for larger values of $n$ will involve many sum of products and will result in $\mathrm{n}^{2}$ equations for a matrix of order n . In this paper, we propose a straightforward method based on multipliers derived from modification of Gaussion elimination algorithm.


KEY WORDS: Lower and Upper Triangular Matrices, Multipliers.

## INTRODUCTION

Let A be a square matrix of order n . An LU factorization or decomposition is a decomposition of the form:
A = LUé é é é é é é é é é é ..

Where L and U are upper and lower triangular matrices (of the same size) respectively (Horn and Johnson, 1985; Kreyszig, 1993; Morris, 1983; Conte, 1965).

The LU factorization is not unique if one only requires that $L$ be lower triangular and $U$ be upper triangular. It is unique if we assign fixed values to the diagonal elements of either L or $U$ (Conte, 1965; Sastry, 1989; Olayi, 2000; Atkinson, 1993).

LU decomposition is used for solving system of linear equations, calculating matrix determinants and inverse.

## THEOREM 1 (EXISTENCE AND UNIQUENESS).

The matrix

$$
\mathrm{A}=\left[\begin{array}{ccc}
a_{11} & a_{12} & \cdots  \tag{2}\\
a_{21} & a_{1 n} \\
a_{22} & a_{22} & a_{n n} \\
a_{n 1} & a_{n 2} & \cdots \\
a_{n n}
\end{array}\right]
$$

admits an LU factorization if and only if all its principal minors are non singular, that is, if
$a_{11} \neq 0 \quad\left|\begin{array}{ll}a_{11} & a_{13} \\ a_{21} & a_{22}\end{array}\right| \neq 0\left|\begin{array}{lll}a_{11} & a_{12} & a_{12} \\ a_{21} & a_{22} & a_{25} \\ a_{81} & a_{82} & a_{82}\end{array}\right| \neq 0$ é é $|A| \neq 0$
(Conte, 1965; Sastry, 1989; Olayi, 2000).

## LU DECOMPOSITION ALGORITHMS

We now outline the various procedures or methods that have hitherto been used to factor a square matrix A into a product of $L$ and $U$ matrices. We assume in all the methods that no interchanges will be necessary. The methods we are going to examine involve writing the general forms of $L$ and $U$ and the unknown elements of $L$ and $U$ are then found by equating corresponding entries in A and LU in a systematic way.
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## DOOLITTLE ALGORITHM

In this algorithm, the lower triangular matrix has all diagonal elements equal to 1, whereas the upper triangular matrix $U$ is of the general form. Thus, the elements of the matrices $L=\left(l_{i j}\right)$ [ with main diagonal1, é, 1] and $U=\left(u_{i j}\right)$ in this method are computed from (Schied, 1988):

$$
\begin{array}{ll}
\mathrm{u}_{\mathrm{ij}}=\mathrm{a}_{1 \mathrm{j}} & \mathrm{j}=1,2, \text { é } \mathrm{n} \\
l_{i \mathbf{1}}=\frac{a_{i \mathbf{1}}}{U_{11}}, & \mathrm{i}=2 \text {, é é } \mathrm{n} \\
\mathrm{u}_{\mathrm{ij}}=\mathrm{a}_{\mathrm{ij}}-\sum_{k=1}^{i-1} l_{i k} u_{k j} & \mathrm{j}=\mathrm{i}, \text { é } \mathrm{n} \\
l_{i j}=\frac{a_{i j}-\sum_{k=1}^{i-1} l_{i k} u_{k i}}{u_{i i}} & \mathrm{i}=\mathrm{j}+\mathrm{l}, \text { é } . . \mathrm{n} \tag{4}
\end{array}
$$

## CROUT'S ALGORITHM

In Crout $\hat{S}$ algorithm, the matrix $U$ has all diagonal elements equal to 1 , whereas $L$ has the general diagonal. Hence, the elements of the matrices
$\mathrm{L}=\left({ }^{l_{i j}}\right)$ and $\mathrm{U}=\left(\mathrm{u}_{\mathrm{ij}}\right)$ [with main diagonal 1, é , 1] are computed from:

$$
\begin{array}{ll}
l_{i \mathbf{1}}=\mathrm{a}_{\mathrm{i1}} & \mathrm{i}=1,2 \text {, é } \mathrm{n} \\
\mathrm{a}_{1 \mathrm{ij}} \\
\mathrm{u}_{11} & \mathrm{j}=2 \text {, é } \mathrm{n} \\
l_{i j}=a_{i j}-\sum_{k=1}^{i-1} l_{i k} u_{k j} & \mathrm{i}=\mathrm{j}, \text {, é } \mathrm{n}  \tag{5}\\
u_{i j}=a_{i j}-\frac{\sum_{k=1}^{i-1} l_{i k} u_{k j}}{l_{i i}} & \mathrm{j}=\mathrm{i}+1, \text { é } \mathrm{n}
\end{array}
$$

## CHOLESKY'S ALGORITHM

For a symmetric positive definite matrix $A\left(A=A^{\top}, x^{\top} A x>0 \forall x \neq 0\right)$. We can choose $U=L^{\top}$, thus $u_{i j}=l_{j i}$ and (4) are simplified to (Kreyszig, 1993)

$$
\left.\begin{array}{ll}
l_{11}=\sqrt{a_{11}} & \\
l_{i i}=\sqrt{a_{11}-\sum_{k=1}^{i-1} l_{i k}^{2}} & \mathrm{i}=2, \text { é .n } \\
l_{i 1}=\frac{a_{i 1}}{l_{11}} &  \tag{6}\\
l_{i j}=\frac{1}{l_{j j}}\left(a_{i j}-\sum_{k=1}^{j-1} l_{i k} l_{k j}\right) & \mathrm{i}=\mathrm{j}+1 \text {,é é n }
\end{array}\right\}
$$

## FACTORIZATION WITH MULTIPIERS

Given an nxn matrix,

$$
\mathrm{A}=\mathrm{a}_{\mathrm{ij}}=\left[\begin{array}{ccc}
a_{11} & a_{12} \cdots & a_{1 n}  \tag{7}\\
a_{21} & a_{22} & a_{2 n} \\
\vdots & \vdots & \vdots \\
a_{n 1} & a_{n 2} \cdots & a_{n n}
\end{array}\right] \quad \text { é é é é é .. }
$$

We want to factor $A$ into the form, $A=L U$

$$
\begin{align*}
\text { With } U & =\left[\begin{array}{cccc}
u_{11} & u_{12} & u_{12} \ldots & u_{n n} \\
u_{21} & u_{22} & u_{22} \ldots & u_{m n} \\
0 & 0 & u_{32} \ldots & u_{m n} \\
0 & 0 & 0 \ldots & u_{m n}
\end{array}\right] \text { é é é é é é .. }  \tag{8}\\
\text { And } \mathrm{L}= & {\left[\begin{array}{cccc}
1 & 0 & 0 \ldots . & 0 \\
l_{21} & 1 & 0 \ldots . & 0 \\
l_{31} & l_{32} & 1 . . & 0 \\
l_{n 1} & l_{n 2} & l_{n 2} \ldots & 1
\end{array}\right]_{\text {é é é é é é é }} } \tag{9}
\end{align*}
$$

Recall the Gaussian elimination algorithm that for a matrix of order $n$, the elimination is performed in ( $n-1$ ) steps, $K=1$,2é .. $n-1$. In step $K$, the elements $a_{i j}^{(k)}$ with $\mathrm{i}, \mathrm{j}>\mathrm{k}$ are transformed according to (Dahlquist and Bjorck; 1974):
$m_{i k}=\frac{a_{i k}(k)}{a_{k k}(k)}$

$a_{i j}^{(k+1)}=a_{i j}^{(k)}-m_{i k} a_{k j}^{(k)^{[-1}}$ é é é é é $\ldots \ldots \ldots$.
i=k+1, $k+2$,é.$n ; \quad j=i, i+1$,é é $n$
Where $\mathrm{m}_{\mathrm{ik}}$ is called the multiplier.
It has been shown by Dahlquist \& Bjorck (1974), Scheid(1988) and Matthews(1987) that the elements in $L$ are the multipliers and the matrix $U$ the final triangular matrix obtained by Gaussian elimination.

Hence, we can say that:
$\mathrm{m}_{\mathrm{ik}}=l_{i k}$
(10) and (11) can now be written as:
$l_{i k}=\frac{\left.a_{i k}{ }^{k}\right)}{a_{k k}{ }^{(k)}}$ é é é é é é é é é é é é é
$a_{i j}^{(k+1)}=a_{i j}^{(k)}-l_{i k} a_{k j}^{(k)}$ é é é é é é é é é
Also, observe that after triangularisation, (7) will take the form:
$\mathrm{U}=\left[\begin{array}{ccccc}a_{11}^{(1)} & a_{12}^{(1)} & a_{12}^{(1)} & \ldots & a_{1 n}^{(1)} \\ 0 & a_{22}^{(2)} & a_{25}^{(2)} & \ldots & a_{2 n}^{(2)} \\ 0 & 0 & a_{3 n}^{(3)} & a_{5 n}^{(3)} \\ 0 & 0 & 0 \ldots & a_{m n}^{(0)}\end{array}\right] \quad$ é é é é.

So, we can let $A=a_{i j}$ in (7) equals $a_{i j}{ }^{(1)}$,
That is, let $A=a_{i j}=a_{i j}^{(1)} \quad$ é é é é é

Comparing (8) with (14), we can say that,
$a_{i j}{ }^{(1)}=u_{i j}, \quad j=1$ to $n \quad$ é é é é é
we already know that,
$l_{\mathrm{ii}}=1, \mathrm{i}=1$ to $\mathrm{n} \quad$ é é é é é é

Instead of writing $\mathrm{i}=\mathrm{k}+1, \mathrm{k}+2$, é. $\mathrm{n} ; \mathrm{j}=\mathrm{i}, \mathrm{i}+1$, é n ; we can write:
$\mathrm{i}=2$ to n for (12), since for $k=1$, this transformation begins from row 2 and $\mathrm{i}, \mathrm{j}=2$ to n for (13) since for $k=1$, it begins from row 2 column 2.
Also comparing (8) with (14), we can say that:
$a_{i j}^{(i)}=u_{i j}, i=2$,é.$n$
é é é é é é

Combining (15), (16), (17), (12), (13) and (18) we now write an algorithm for factoring A into LU:

$$
\begin{aligned}
& \operatorname{Let} A=a_{i j}=a_{i j}{ }^{(1)} \\
& a_{1 j}{ }^{(1)}=u_{i j}, j=1 \text { to } n \\
& l_{\mathrm{ii}}=1, \mathrm{i}=1 \text { to } \mathrm{n} \\
& \text { For } \mathrm{k}=1,2 \text { to } \mathrm{n}-1 \\
& l_{i k}=\frac{a_{i k}^{(k)}}{a_{k k^{(k)}}^{(k)}} \quad \text { i>k, i }=2 \text { to } \mathrm{n} \\
& a_{i j}^{(k+1)}=a_{i j}^{(k)}-l_{i k} a_{k j}^{(k)}{ }_{i, j}^{(k)}+i, j=2 \text { to } n: \\
& a_{i j}^{(i)}=u_{i j} i, j=2 \text {,é } n \\
& U=\left(u_{i j}\right) 1 \text { Öi } i, j \text { OOn and } L=\left(L_{i j}\right) \text { 1Òi, } j \text { Òn }
\end{aligned}
$$

APPLICATION (Stroud, 1996)
We want to decompose
$A=\left[\begin{array}{ccc}3 & 2 & -1 \\ 2 & -1 & 2 \\ 1 & -3 & -4\end{array}\right]$ into $A=L U$,
Which we know the result to be:


METHOD 1: USING MULTIPLIERS
$a_{11}{ }^{(1)}=3, a_{12}{ }^{(1)}=2, a_{13}{ }^{(1)}=-1, a_{21}{ }^{(1)}=2, a_{22}{ }^{(1)}=-1, a_{23}{ }^{(1)}=2$
$a_{31}{ }^{(1)}=1, a_{32}{ }^{(1)}=-3, a_{33}{ }^{(1)}=-4$.
$a_{1 j}{ }^{(1)}=u_{1 j}, j=1$ to $n \Rightarrow$
$a_{11}{ }^{(1)}=u_{11}=3, a_{12}{ }^{(1)}=u_{12}=2, a_{13}{ }^{(1)}=u_{13}=-1$
$I_{11}=1, \mathrm{i}=1$ to $\mathrm{n} \Rightarrow$
$I_{11}=I_{22}=l_{33}=1$
For $k=1$ to $n-1$, we have:
$\mathrm{K}=1, \mathrm{i}=2, \Rightarrow I_{21}=2 / 3$
$K=1, i=3, \Rightarrow I_{31}=1 / 3$
$K=1, i=2, j=2 \Rightarrow a_{22}{ }^{(2)}=-7 / 3$
$\mathrm{K}=1, \mathrm{i}=2, \mathrm{j}=3 \Rightarrow \mathrm{a}_{23}{ }^{(2)}=8 / 3$
$K=1, i=3, j=2 \Rightarrow a_{32}{ }^{(2)}=-11 / 3$
$K=1, i=3 j=3, \Rightarrow a_{33}{ }^{(2)}=-11 / 3$
$\mathrm{K}=2, \mathrm{i}=3, \Rightarrow \mathrm{I}_{32}=11 / 7$
$\mathrm{K}=2, \mathrm{i}=3 \mathrm{j}=3, \Rightarrow \mathrm{a}_{33}{ }^{(3)}=-55 / 7$

Thus,
$a_{22}{ }^{(2)}=U_{22}=-7 / 3, a_{23}{ }^{(2)}=U_{23}=8 / 3, \quad a_{22}{ }^{(3)}=U_{33}=-55 / 7$,
$\therefore L=\left[\begin{array}{ccc}\frac{1}{2} & 0 & 0 \\ \frac{1}{3} & 1 & 0 \\ \frac{1}{3} & \frac{11}{7} & 1\end{array}\right] \quad U=\left[\begin{array}{ccc}3 & 2 & -1 \\ 0 & -\frac{7}{3} & \frac{0}{3} \\ 1 & 0 & -\frac{55}{7}\end{array}\right]$
METHOD 2: USING DOOLITTLE ALGORITHM
For the purpose of our comparison, we shall use Doolittle algorithm.
We already know that, Doolittle algorithm(4) is obtained by writing the general forms of $L$ and $U$, where $L$ has all the diagonal elements equal to, whereas the upper triangular matrix $U$ is of the general form and the unknown elements of $L$ and $U$ are then found by equating corresponding entries in $A$ and $L U$ in a systematic way. Thus, for:

A

$$
=\left[\begin{array}{ccc}
3 & 2 & -1 \\
2 & -1 & 2 \\
1 & 3 & -4
\end{array}\right]
$$

Let $I_{11}=\quad I_{22}=\quad I_{33}=1$
$\mathrm{LU} \quad=\left[\begin{array}{ccc}1 & 0 & 0 \\ l_{21} & 1 & 0 \\ l_{31} & l_{32} & 1\end{array}\right]\left[\begin{array}{ccc}U_{11} & U_{12} & U_{12} \\ 0 & U_{22} & U_{23} \\ 0 & 0 & U_{32}\end{array}\right]$


But $A=L U \Rightarrow U_{11}=3, U_{12}=2, \quad U_{13}=-1$
$l_{21} u_{11}=2, \Rightarrow 3 l_{21}=2, \Rightarrow l_{21}=2 / 3$
$l_{31} u_{11}=1 \Rightarrow 3 l_{31}=1, \Rightarrow l_{31}=1 / 3$
$l_{21} u_{12}+u_{22}=-1 \quad \Rightarrow 4 / 3+u_{22}=-1$
$\Rightarrow u_{22}=-1-4 / 3=-7 / 3$
$l_{21} u_{13}+u_{2 a}=2, \quad \Rightarrow-2 / 3+\mathrm{U}_{23}=2$
$\Rightarrow u_{23}=2+2 / 3=8 / 3$
$l_{31} u_{12}+l_{32} u_{22}^{\text {TI }}=-3, \quad \Rightarrow 1 / 3^{(2)}+l_{32}^{\left(\frac{7}{3}\right)}=3$
$\Rightarrow l_{32}=11 / 7$
$l_{31} U_{13}+l_{32} U_{23}+U_{3 n}=-4$
$1 / 3(-1)+11 / 7(8 / 3)+U_{35}=-4$
$-1 / 3+88 / 21+U_{33}=-4$
$U_{\text {sa }}=-4$ ї $81 / 21=-165 / 21=-55 / 7$
Thus $A=L U=\left[\begin{array}{ccc}\frac{1}{2} & 0 & 0 \\ \frac{1}{3} & 1 & 0 \\ \frac{1}{3} & \frac{11}{7} & 1\end{array}\right]\left[\begin{array}{ccc}3 & 2 & -1 \\ 0 & -\frac{7}{3} & \frac{8}{3} \\ 0 & 0 & -\frac{55}{7}\end{array}\right]$

## CONCLUSION

We have modified the Gaussian elimination algorithm and have developed a straightforward algorithm based on multipliers for factoring an $n \times n$ matrix $A$ into the form $A=L U$, where $L$ are the multipliers with Is on the diagonal and $U$ is the upper triangular matrix. We have also observed that our proposed algorithm does not involve many sums of products as compared to the Dool ittle algorithm.

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